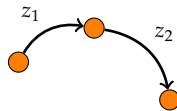


ARTUR EKERT

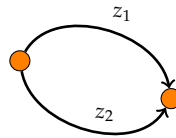
About complex numbers, called probability amplitudes, that, unlike probabilities, can cancel each other out, leading to quantum interference and qualitatively new ways of processing information. About impossible logic operations, such as $\sqrt{\text{NOT}}$, which can be implemented. And about qubits and a single qubit interference.

Quantum mechanics, at least at some instrumental level, can be viewed as a modification of the probability theory. We replace positive numbers (probabilities) with complex numbers z (probability amplitudes) such that the squares of their absolute values, $|z|^2$, are interpreted as probabilities. The rules for combining amplitudes are very reminiscent of the rules for combining probabilities.



$$z = z_1 z_2$$

Whenever something can happen in a sequence of independent steps we multiply the amplitudes of each step.

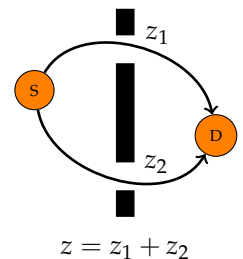


$$z = z_1 + z_2$$

Whenever something can happen in several alternative ways we add amplitudes for each way considered separately.

That's it! The two rules are basically all you need to manipulate amplitudes in any, no matter how complicated, physical process (we will amend the two rules later on when we touch upon the particle statistics). The two simple rules are universal and apply to any physical system, to the big and to the small, from elementary particles through atoms and molecules to white dwarfs stars. They also apply to information for information is physical; it is both represented and processed by physical means. Here we will use the two rules to explain a distinctive power of quantum information processing.

1.1. Quantum Interference. In order to see the need for quantum theory let us consider a simple experiment in which probability theory fails to give the right predictions. In a double slit experiment a particle emitted from a source S can reach detector D by taking either an upper or a lower slit, with amplitudes z_1 and z_2 respectively. We may say that the upper slit is taken with probability $p_1 = |z_1|^2$ and the lower slit with probability $p_2 = |z_2|^2$. These are two mutually exclusive events. With the two slits open, probability theory declares, the particle will reach the detector with probability $p_1 + p_2 = |z_1|^2 + |z_2|^2$. Wrong! The particle will reach the detector with probability

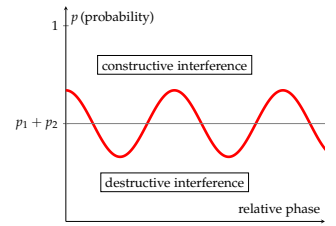


$$\begin{aligned}
 p = |z|^2 = |z_1 + z_2|^2 &= |z_1|^2 + |z_2|^2 + z_1^* z_2 + z_1 z_2^*, \\
 &= p_1 + p_2 + |z_1| |z_2| (e^{-i(\phi_2 - \phi_1)} + e^{i(\phi_2 - \phi_1)}), \\
 &= p_1 + p_2 + 2\sqrt{p_1 p_2} \cos(\phi_2 - \phi_1), \\
 &= p_1 + p_2 + \text{interference terms}, \tag{1}
 \end{aligned}$$

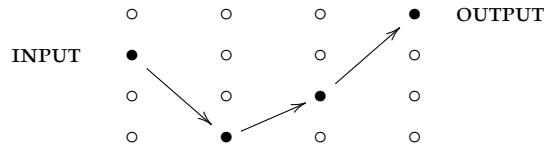
where we have expressed the amplitudes in their polar form $z_1 = |z_1|e^{i\phi_1}$ and $z_2 = |z_2|e^{i\phi_2}$. The appearance of the interference terms marks the departure from the classical theory of probability. The probability of any two seemingly mutually exclusive events is the sum of the probabilities of the individual events, $p_1 + p_2$,

Add amplitudes, not probabilities. It seems that nobody told nature about the Kolmogorov additivity axiom.

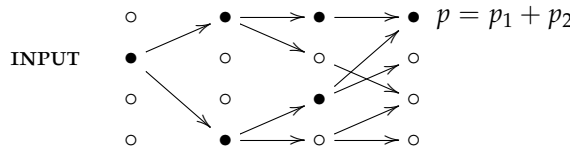
modified by the interference term, $2\sqrt{p_1 p_2} \cos(\phi_2 - \phi_1)$. Depending on the relative phase $\phi_2 - \phi_1$, the interference term can be either negative (destructive interference) or positive (constructive interference), leading to either suppression or enhancement of the total probability p . If we can control relative phases we can take advantage of this effect. For example, quantum computation can be viewed as a complex multiparticle quantum interference, involving many computational paths, which is designed to enhance probabilities of correct outputs and to suppress probabilities of the wrong ones.



1.2. Deterministic, probabilistic and quantum. Think about computation as a physical process that evolves a prescribed initial configuration of a computing machine, called input, into some final configuration, called output. The diagram below shows three consecutive computational steps performed on four distinct configurations.

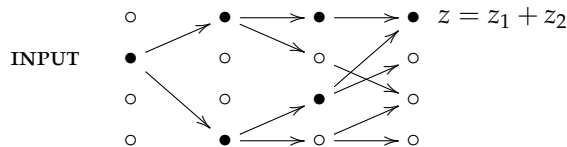


Such a computation does not have to be deterministic. We can augment a computing machine by allowing it "to toss an unbiased coin" and to choose its steps randomly. It can then be viewed as a directed, tree-like graph where each node corresponds to a configuration of the machine, and each edge represents one step of the computation.



Randomised algorithms can be more powerful than deterministic ones

The computation starts from the root node representing the initial configuration and it subsequently branches into other nodes representing configurations reachable with non-zero probability from the initial configuration. The probability of a particular final configuration being reached is equal to the sum of the probabilities along all mutually exclusive paths which connect the initial configuration with that particular configuration. The snag is that it does not always work this way. There are many physical phenomena, usually involving atoms or photons, which blatantly ignore the additivity axiom. Thus, in order to make statistical predictions that agree with experiments, probability theory had to be modified. A quantum computation can be represented by a graph similar to that of a probabilistic computation.



Quantum algorithms can be more powerful than randomised ones. Quantum interference is an extra tool to our disposal.

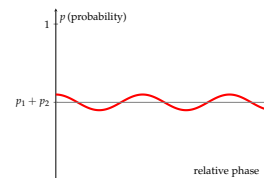
Following the rules of quantum theory we associate with each edge in the graph the probability *amplitude* that the computation follows that edge. The probability amplitude of a particular final configuration being reached is equal to the sum of the amplitudes along all mutually exclusive paths which connect the initial configuration with that particular configuration. The resulting probability, as we have just seen, reads $p_1 + p_2 + 2\sqrt{p_1 p_2} \cos(\phi_2 - \phi_1)$. The basic idea behind quantum algorithms is to use quantum interference to amplify the correct outputs and to suppress all other outcomes of computations. It requires a significant control over relative phases in a quantum computer for the phase settings control the interference and hence the computation.

1.3. What is wrong with the additivity axiom? You may be wondering what has happened to the axiom of additivity in probability theory, which says that if E_1 and E_2 are mutually exclusive events then the probability of the event (E_1 or E_2) is the sum of the probabilities of the constituent events, E_1, E_2 . So what is wrong with the additivity axiom? One thing that is wrong is the assumption that the processes of taking the upper or the lower slit are mutually exclusive. In reality, the two transitions *both occur*, simultaneously. We cannot learn about this fact from probability theory or any other a priori mathematical construct. We learn it from the best physical theory available at present, namely quantum theory. This knowledge was created as the result of conjectures, experimentation, and refutations.

Here I shamelessly reveal my philosophical leaning: Karl Popper and his *Conjectures and Refutations: The Growth of Scientific Knowledge*

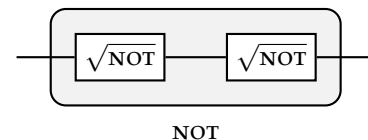
1.4. Relative Phase. Note that the important quantity here is the relative phase $\phi_1 - \phi_2$ rather than the absolute values ϕ_1 and ϕ_2 . This observation is not trivial at all. In simplistic terms - if a particle reacts only to the difference of the two phases, each pertaining to a separate path, then it must have, somehow, experienced the two paths, right? Thus we cannot say that the particle has travelled either through the upper or the lower slit, it has travelled through *both*. In the same way quantum computers follow, in some tangible way, all computational paths simultaneously, producing answers that depend on all these alternative calculations. Weird but this is how it is!

1.5. Decoherence. If this is how it is, so why we do not see quantum interference on a daily basis? This is because phases of probability amplitudes tend to be very fragile and may fluctuate rapidly due to spurious interactions with the environment. This has influence on the interference term; it may average to zero and we recover the classical addition of probabilities. This phenomenon is known as *decoherence*. We will discuss the origin of decoherence later on, for now let me only mention that it is very conspicuous in physical systems made out of many interacting components and it is chiefly responsible for our classical description of the world – without interference terms we may as well add probabilities instead of amplitudes. Decoherence, as you can imagine, is a serious impediment to building quantum computers; it deprives us of the power of quantum interference. This said, it is not all doom and gloom, there are clever ways around decoherence and later on we will touch upon quantum error corrections and quantum fault-tolerant computations.



Decoherence suppresses quantum interference.

1.6. Impossible logic. Now that we have poked our heads into the quantum world, let us see how quantum interference challenges our good old conventional logic and leads to qualitatively different computations. Consider a following task: design a logic gate that operates on a single bit and such that when it is followed by another, identical, logic gate the output is always the negation of the input. Let us call this logic gate the square root of NOT ($\sqrt{\text{NOT}}$). A simple check – such as an attempt to construct a truth table – should persuade you that there is no such operation in logic. It may seem reasonable to argue that since there is no such operation in logic, $\sqrt{\text{NOT}}$ is impossible. But it does exist! Experimental physicists routinely construct such “impossible” gates in their laboratories. In fact the $\sqrt{\text{NOT}}$ can be as simple as a beam-splitter.



1.7. Physics against logic. A symmetric beam-splitter is a cube of glass which reflects half the light that impinges upon it, while allowing the remaining half to pass through unaffected. For our purposes it can be viewed as a device which has two input and two output ports which we label as $|0\rangle$ and $|1\rangle$. When we aim a single photon at such a beam-splitter using one of the input ports, we notice that the photon doesn't split in two: we can place photo-detectors wherever we like in the apparatus, fire in a photon, and verify that if any of the photo-detectors registers a hit, none of the others do. In particular, if we place a photo-detector behind the beam-splitter in each of the two possible exit beams, the photon is detected with

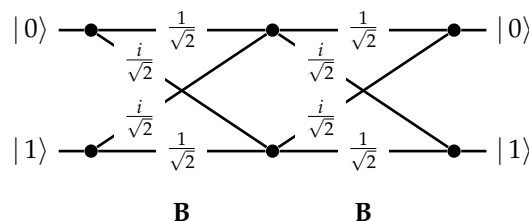
equal probability at either detector, no matter whether the photon was initially fired from input port $|0\rangle$ or $|1\rangle$. It may seem obvious that at the very least, the photon is *either* in the transmitted beam $|0\rangle$ or in the reflected beam $|1\rangle$ during any one run of this experiment. Thus we may be tempted to think of the beam-splitter as a random binary switch which, with equal probability, transforms any binary input into one of the two possible outputs. However, that is not necessarily the case. Let us introduce a second beam-splitter and place two normal mirrors so that both paths intersect at the second beam-splitter (see diagrams in the margin).

Now, the axiom of additivity in probability theory, says that whenever something can happen in several alternative ways we add probabilities for each way considered separately. We might argue that a photon fired into the input port $|0\rangle$ can reach the detector 0 in two *mutually exclusive* ways: either by two consecutive reflections or by two consecutive transmissions. Each reflection happens with probability $1/2$ and each transmission happens with probability $1/2$ thus the total probability of reaching detector 0 is a sum of the probability of the two consecutive reflections ($1/2 \times 1/2 = 1/4$) and the probability of the two consecutive transmissions ($1/2 \times 1/2 = 1/4$) which gives probability $1/2$. This makes perfect sense – a random switch followed by a random switch should give nothing else but a random switch. However, if we set up such an experiment, that is not what happens! When the optical paths between the two beam-splitters are the same, the photon fired from input port $|0\rangle$ *always* strikes detector 1 and *never* detector 0 (and the photon fired from input port $|1\rangle$ *always* strikes detector 0 and *never* detector 1). Thus a beam-splitter acts as the square root of NOT gate. What is wrong with our reasoning here? Why does probability theory fail to predict the outcome of this simple experiment?

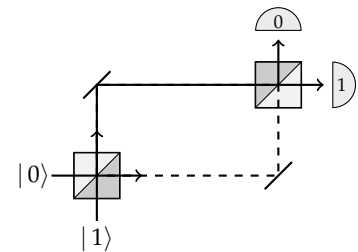
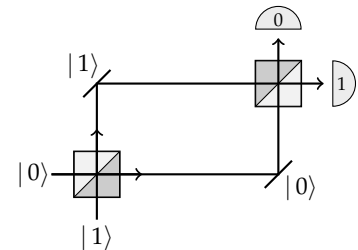
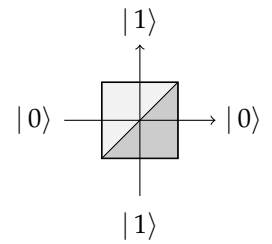
1.8. Enter quantum theory. A beamsplitter is a quantum device thus we have to multiply and add amplitudes not probabilities. The action of the beamsplitter – in fact, the action of any quantum device – can be described by tabulating the amplitudes of transitions between its input and output ports.

$$B = \begin{bmatrix} B_{00} & B_{01} \\ B_{10} & B_{11} \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{i}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} & \frac{1}{\sqrt{2}} \end{bmatrix}.$$

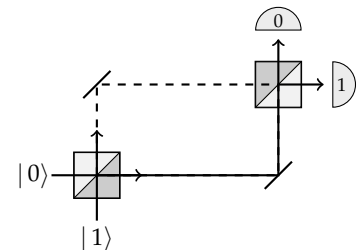
The matrix element B_{lk} , where $k, l = 0, 1$, represents the amplitude of transition from input $|k\rangle$ to output $|l\rangle$ (watch the order of indices). Each reflection (entries B_{01} and B_{10}) happens with amplitude $i/\sqrt{2}$ and each transmission (entries B_{00} and B_{11}) happens with amplitude $1/\sqrt{2}$. Thus the total amplitude that a photon fired from input port $|0\rangle$ will reach detector 0 is the sum of the amplitude of the two consecutive reflections ($i/\sqrt{2} \times i/\sqrt{2} = -1/2$) and the amplitude of the two consecutive transmissions ($1/\sqrt{2} \times 1/\sqrt{2} = 1/2$) which gives the total amplitude 0. The resulting probability is then zero. Unlike probabilities, amplitudes can cancel out each other out. We can now go on and calculate the amplitude that the photon will reach detector 1. In this case we will get i , which gives probability 1. We can then switch to input $|1\rangle$ and repeat our calculations. All possible paths and associated amplitudes are shown in the diagram below.



However, instead of going through all the paths in this diagram and linking specific inputs to specific outputs, we can simply multiply the transition matrices,



Two consecutive reflections give amplitude $\frac{i}{\sqrt{2}} \frac{i}{\sqrt{2}} = -\frac{1}{2}$



Two consecutive transmissions give amplitude $\frac{1}{\sqrt{2}} \frac{1}{\sqrt{2}} = \frac{1}{2}$

There is no reason why probability theory or any other a priori mathematical construct should make any meaningful statements about outcomes of physical experiments.

LOGICAL NOT

$$X = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

— [X] —

$\sqrt{\text{NOT}}$

$$B = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & i \\ i & 1 \end{bmatrix}$$

— [B] —

$$BB = \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{i}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} & \frac{1}{\sqrt{2}} \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{i}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} & \frac{1}{\sqrt{2}} \end{bmatrix} = \begin{bmatrix} 0 & i \\ i & 0 \end{bmatrix} = iX.$$

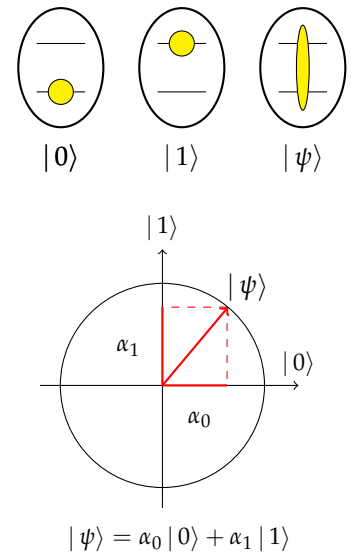
As you can see, the matrix multiplication in one swoop takes care of multiplication and addition of amplitudes corresponding to different alternatives. Quantum theory explains the behaviour of $\sqrt{\text{NOT}}$ and correctly predicts the probabilities of all the possible outputs. Hence, reassured by the physical experiments that corroborate this theory, logicians are now entitled to propose a new logical operation $\sqrt{\text{NOT}}$. Why? Because a faithful physical model for it exists in nature!

Gate B is not the only quantum operation that effects the square root of NOT. Can you find any other? There are infinitely many of them.

1.9. Quantum bits called qubits. Needless to say, a symmetric beam-splitter is just one way of implementing the square root of NOT – there are many others. For example, instead of choosing between two different beams of light a bit can also be encoded by choosing between two different states of an atom, e.g. state $|0\rangle$ may be chosen to be the lowest energy state, the ground state, and state $|1\rangle$ a higher energy state, the excited state. Pulses of light of appropriate frequency, duration and intensity can take the atom back and forth between the states $|0\rangle$ and $|1\rangle$ (implementing logical NOT). Some other pulses, say, half the duration or intensity, which implement $\sqrt{\text{NOT}}$, will take the atom into states that have no classical analogue; they are called *coherent superpositions* of $|0\rangle$ and $|1\rangle$. Any object in which such states can be reliably prepared, manipulated and measured is called a quantum bit or a *qubit*. What are these qualitatively new states? Technically, a coherent superposition means the qubit is in state $|0\rangle$ with some amplitude α_0 and in state $|1\rangle$ with some other amplitude α_1 . This is conveniently represented by a state vector

$$|\psi\rangle = \alpha_0 |0\rangle + \alpha_1 |1\rangle \leftrightarrow \begin{bmatrix} \alpha_0 \\ \alpha_1 \end{bmatrix}. \tag{2}$$

We square the absolute values of amplitudes to get probabilities; if we choose to measure the bit value (say, energy of the atom) in such a superposition we will find either 0 (the energy of the ground state) or 1 (the energy of the excited state), with probabilities $|\alpha_0|^2$ and $|\alpha_1|^2$, respectively. As no other states are involved these are the only two possible and mutually exclusive outcomes of the measurement hence we require that $|\alpha_0|^2 + |\alpha_1|^2 = 1$. By analogy with the elementary geometry we can view this expression as the square of the length of the state vector and declare that all state vectors are of unit length. In general, any isolated quantum object with n perfectly distinguishable states, or configurations, can be completely described by a unit vector in an n -dimensional complex vector space, $|\psi\rangle = \sum_k \alpha_k |k\rangle$, where $k = 0, 1, \dots, n - 1$ and α_k is the amplitude that the system is in state $|k\rangle$. Any physically admissible operation on this object is then completely described by an $n \times n$ matrix of complex numbers. For example, matrix B describes an admissible physical operation on a qubit. As these matrices map state vectors into state vectors, that is, unit vectors into unit vectors, they must be isometries (operations that preserve the length) which implies that they must be unitary.



Here we are talking about qubits but all constructions and formulae can be easily generalised to arbitrary finite dimensional spaces.

1.10. Unitary evolutions. Any quantum evolution U on an isolated system is represented by a linear operator which sends state $|k\rangle$ to state $|l\rangle$ ($k, l = 0, 1$) but only with some *amplitude* U_{lk} (watch the order of the indices). We write this as

$$|k\rangle \rightarrow \sum_l U_{lk} |l\rangle. \tag{3}$$

The operator is succinctly specified by the matrix U_{lk} , which tabulates all possible transition amplitudes between states $|0\rangle$ and $|1\rangle$. As state vectors are unit vectors, all admissible quantum evolutions must be represented by isometries. In Euclidean space the relevant isometries (rotations, reflections, and their combinations) are described by orthogonal matrices. Here we are dealing with vectors which have complex components; if we multiply these components by phase factors, such as

$e^{i\varphi}$, the length does not change. This leads to a larger class of isometries described by unitary matrices. Recall that matrix U is called unitary if

$$U^\dagger U = U U^\dagger = \mathbb{1}, \tag{4}$$

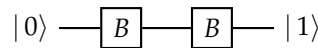
where the *adjoint* or *Hermitian conjugate* U^\dagger of any matrix U with complex entries U_{ij} is obtained by taking the complex conjugate of every element in the matrix and then interchanging rows and columns ($U_{kl}^\dagger = U_{lk}^*$). If you prefer to write the matrices explicitly, in terms of their components, then the unitarity condition above can be expressed as $\sum_l U_{lm}^* U_{lk} = \sum_l U_{ml} U_{kl}^* = \delta_{mk}$. Operator U sends state $|\Psi\rangle$, with components α_k , into state $|\Psi'\rangle = U|\Psi\rangle$, with components $\alpha'_l = \sum_k U_{lk} \alpha_k$,

$$|\Psi\rangle = \sum_k \alpha_k |k\rangle \rightarrow \sum_k \alpha_k \left(\sum_l U_{lk} |l\rangle \right) = \sum_l \left(\sum_k U_{lk} \alpha_k \right) |l\rangle = \sum_l \alpha'_l |l\rangle = |\Psi'\rangle. \tag{5}$$

Isometries include both unitary and anti-unitary operators, such as time reversal, about which later.

Here δ_{mk} , known as the "Kronecker delta", is a symbol that is defined to be zero for $k \neq m$ and to be one for $k = m$

1.11. Quantum circuits. Atoms, trapped ions, molecules, nuclear spins and many other quantum objects can be prepared not only in two distinct states, labelled as $|0\rangle$ and $|1\rangle$, but also in any other quantum state described by a state vector. There is no need to learn about physics behind these diverse technologies if all you want is to understand the foundations of quantum computation. For example, to understand the square root of NOT you may conveniently forget about any specific experimental realisation of this gate and draw a diagram which represents them all,



This quantum network, or quantum circuit, diagram should be read from left to right. The horizontal line represents a qubit that is inertly carried from one quantum operation to another. We often call it a quantum wire. The wire may describe translation in space, e.g. atoms traveling through cavities, or translation in time, e.g. a sequence of operations performed on a trapped ion. If we want to signify that a particular unitary evolution is to be enacted on our qubit, then we put a box with a symbol describing this unitary operation along the quantum wire. You can step through the diagram and follow the evolution of the state vector

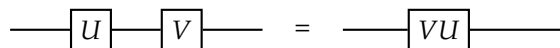
$$|0\rangle \xrightarrow{B} \frac{1}{\sqrt{2}} (|0\rangle + i|1\rangle) \xrightarrow{B} i|1\rangle. \tag{6}$$

State $i|1\rangle$ is the same as state $|1\rangle$, the **global** phase factor can be ignored. In general, state $e^{i\phi}|\psi\rangle$ and state $|\psi\rangle$ represent the same physical situation. It is a relative phase within a superposition that does matter.

If you prefer to work with column vectors and matrices, you can write the two consecutive application of B to state $|0\rangle$ as

$$\begin{bmatrix} 0 \\ i \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{i}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} & \frac{1}{\sqrt{2}} \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{bmatrix} \leftarrow \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{i}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} & \frac{1}{\sqrt{2}} \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix}. \tag{7}$$

Following a well established convention the formulae above should be read from right to left. Confused? Well, you are not the only one. Just remember that circuits diagrams are read from left to right and vector and matrix operations go from right to left. A circuit composed of two gates, say U followed by V , is equivalent to a circuit composed of one gate described by the matrix product VU (note the order in which we multiply the matrices),



Indeed, if U and V are defined by

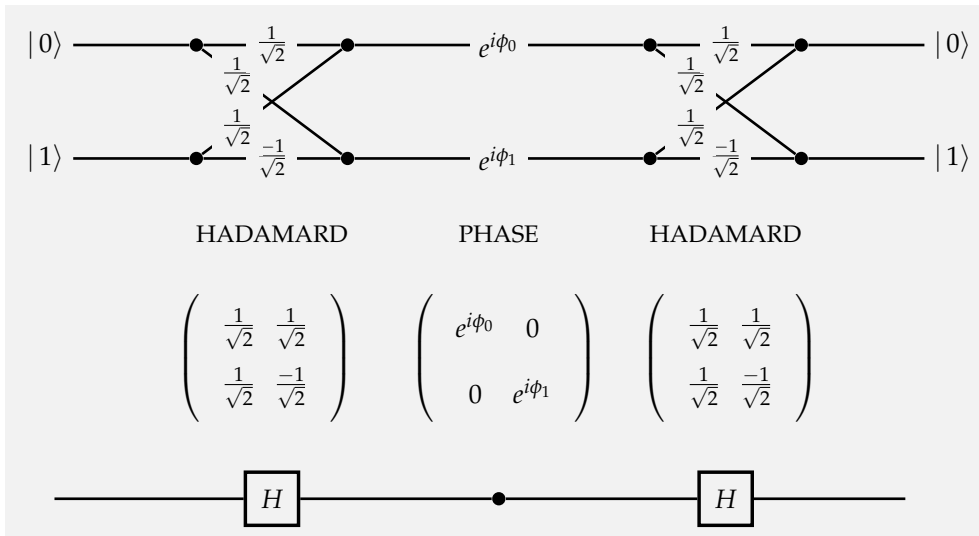
$$|k\rangle \rightarrow \sum_l U_{lk} |l\rangle, \quad |l\rangle \rightarrow \sum_m V_{ml} |m\rangle, \tag{8}$$

then

$$|k\rangle \rightarrow \sum_l U_{lk} \left(\sum_m V_{ml} |m\rangle \right) = \sum_m \left(\sum_l V_{ml} U_{lk} \right) |m\rangle. \tag{9}$$

If you want to hone your quantum intuition think about it this way. The amplitude that input $|k\rangle$ evolves to $|m\rangle$ via a specific intermediate state $|l\rangle$ is given by $V_{ml}U_{lk}$ (evolutions are independent so the amplitudes are multiplied). This done we have to sum over all possible values of l (the transition can occur in several mutually exclusive ways so the amplitudes are added) to obtain $\sum_l V_{ml}U_{lk}$. As you we have already observed, the matrix multiplication in one swoop takes care of multiplication and addition of amplitudes corresponding to different alternatives.

1.12. Single qubit interference. Let me now describe what is probably the most important sequence of operations performed on a single qubit, namely a generic single qubit interference. It is typically constructed as a sequence of three elementary operations: the Hadamard gate, followed by a phase shift gate, and followed by the Hadamard gate. We represent it graphically as



The top diagram helps to visualise different paths connecting the two input and output states of a qubit. Right below you find the three basic operations defined by their corresponding matrices. The bottom diagram is a quantum circuit diagram, our preferred way of representing quantum evolutions. Let us start with brute-force approach and multiply the three matrices. The resulting matrix describes the action of the whole circuit; it gives the transition amplitudes between states $|0\rangle$ and $|1\rangle$ at the input and the output.

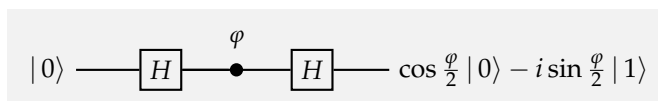
$$\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} e^{i\phi_0} & 0 \\ 0 & e^{i\phi_1} \end{bmatrix} \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} = e^{i\frac{\phi_0+\phi_1}{2}} \begin{bmatrix} \cos \varphi/2 & -i \sin \varphi/2 \\ -i \sin \varphi/2 & \cos \varphi/2 \end{bmatrix},$$

where $\varphi = \varphi_1 - \varphi_0$. Multiplying a state vector by a scalar is physically inconsequential thus we ignore the global phase factor $e^{i(\varphi_0+\varphi_1)/2}$. What really matters is the phase difference φ . Thus the action of the circuit is described by the map

$$|0\rangle \rightarrow \cos \frac{\varphi}{2} |0\rangle - i \sin \frac{\varphi}{2} |1\rangle, \tag{10}$$

$$|1\rangle \rightarrow -i \sin \frac{\varphi}{2} |0\rangle + \cos \frac{\varphi}{2} |1\rangle. \tag{11}$$

Given that our input state is almost always $|0\rangle$ it is sometimes much easier and more instructive to step through the execution of this circuit and follow the evolving state. The interference circuit effects the following sequence of transformations,

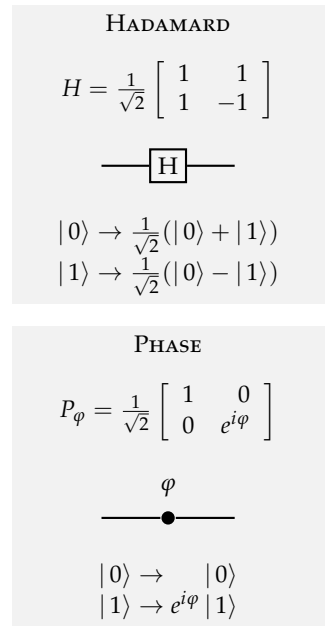


$$|0\rangle \xrightarrow{H} \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \xrightarrow{P_\phi} \frac{1}{\sqrt{2}}(|0\rangle + e^{i\phi}|1\rangle) \xrightarrow{H} \cos\frac{\phi}{2}|0\rangle - i\sin\frac{\phi}{2}|1\rangle. \quad (12)$$

The first Hadamard gate prepares an equally weighted superposition of $|0\rangle$ and $|1\rangle$ and the second one closes the interference by bringing the interfering paths together. The phase shift ϕ effectively controls the evolution and determines the output. The probabilities of finding the qubit in state $|0\rangle$ or $|1\rangle$ at the output are, respectively,

$$\text{Pr}(0) = \cos^2\frac{\phi}{2}, \quad \text{Pr}(1) = \sin^2\frac{\phi}{2}. \quad (13)$$

This simple quantum process contains, in a nutshell, the essential ingredients of quantum computation. The sequence, Hadamard - phase shift - Hadamard, will appear over and over again. It reflects a natural progression of quantum computation: first we prepare different computational paths, then we evaluate a function which effectively introduces phase shifts into different computational paths, then we bring the computational paths together at the output.



STATES AND QUANTUM EVOLUTIONS

There are essentially four components of the mathematical structure of quantum theory. We need to know how to represent (1) states, (2) physically admissible operations or quantum evolutions, (3) measurements, and (4) composite systems. I will describe them in stages, starting with a very simple case of a single qubit, and working towards the most general case. So far we talked about

- Quantum states are described by unit vectors with complex components. Any two state vectors which differ by a multiplicative phase factor $e^{i\phi}$, e.g. $|\Psi\rangle$ and $e^{i\phi}|\Psi\rangle$, represent the same quantum state. Global phase factors can be ignored. In contrast, all relative phase factors are very important; the two vectors $|\Psi\rangle = \alpha_0|0\rangle + \alpha_1|1\rangle$ and $|\Psi\rangle = \alpha_0|0\rangle + e^{i\phi}\alpha_1|1\rangle$ represent two different states.
- States evolve in time. Any admissible physical evolution U of a closed system is represented by a unitary operator.

Measurements and composite systems will feature soon.

NOTES

(1) Back in 1926 Max Born simply postulated the connection between amplitudes and probabilities. However, it is worth pointing out, that he did not get it quite right on his first approach. In the original paper proposing the probability interpretation of the state vector (wavefunction)¹ he wrote:

...If one translates this result into terms of particles only one interpretation is possible. $\Theta_{\eta,\tau,m}(\alpha, \beta, \gamma)$ [the wavefunction for the particular problem he is considering] gives the probability* for the electron arriving from the z direction to be thrown out into the direction designated by the angles α, β, γ ...

* Addition in proof: More careful considerations show that the probability is proportional to the square of the quantity $\Theta_{\eta,\tau,m}(\alpha, \beta, \gamma)$.

¹Max Born, Zur Quantenmechanik der Stoßvorgänge, *Zeitschrift für Physik*, 37, 863–867 (1926).

- (2) All physical evolutions U are represented by operators that map unit state vectors $\sum_k \alpha_k |k\rangle$ into unit state vectors $\sum_l \alpha'_l |l\rangle$, hence

$$1 = \sum_l |\alpha'_l|^2 = \sum_l \left(\sum_m U_{lm} \alpha_m \right) \left(\sum_n U_{ln} \alpha_n \right)^* = \sum_{m,n} \left(\sum_l U_{lm} U_{ln}^* \right) \alpha_m \alpha_n^*. \quad (14)$$

This equality must hold for all admissible values of α_m and α_n^* . This is possible only if $\sum_l U_{lm} U_{ln}^* = \sum_l U_{nl}^\dagger U_{lm} = \delta_{mn}$, i.e. the operators must be unitary.

EXERCISES

- (1) A quantum computer starts calculations in some initial state, then follows n different computational paths which lead to the final output. The computational paths are followed with probability amplitudes $\frac{1}{\sqrt{n}} e^{ik\varphi}$, where φ is a fixed angle $0 < \varphi < 2\pi$ and $k = 0, 1, \dots, n-1$. Show that the probability of generating the output is

$$1 + z + z^2 + \dots + z^n = \frac{1 - z^{n+1}}{1 - z}$$

$$\frac{1}{n} \left| \frac{1 - e^{in\varphi}}{1 - e^{i\varphi}} \right|^2 = \frac{1}{n} \frac{\sin^2(n\frac{\varphi}{2})}{\sin^2(\frac{\varphi}{2})}.$$

for $0 < \varphi < 2\pi$ and 1 for $\varphi = 0$. Plot the probability as a function of φ .

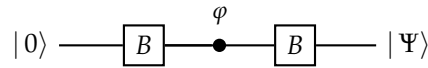
- (2) The four matrices below describe all possible computations on a single classical bit.

$$\begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \quad \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \quad \begin{bmatrix} 1 & 1 \\ 0 & 0 \end{bmatrix} \quad \begin{bmatrix} 0 & 0 \\ 1 & 1 \end{bmatrix}$$

IDENTITY NOT CONST 0 CONST 1

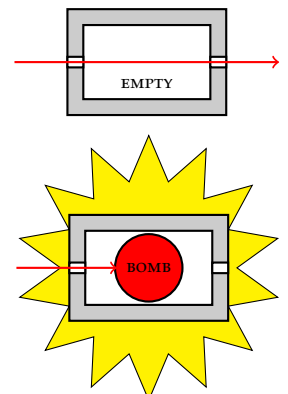
Which of the four operations are also quantum operations?

- (3) A quantum interference experiments is described by the following circuit



where the gate B is a symmetric beamsplitter and the central gate is a phase shift gate which introduces a relative phase ϕ . Step through the execution of this circuit and follow the evolving input state. What is the state of the qubit at the output, here represented as $|\Psi\rangle$? Suppose you can control the input of the circuit and measure the output, but you do not know the phase shift ϕ introduced by the phase gate. You are promised, however, that ϕ is either 0 or π . You can run the circuit only once to find out which of the two phases was chosen?

- (4) **The Quantum Bomb Tester.**² You have been drafted by the government to help in the demining effort in a former war-zone. In particular, retreating forces have left very sensitive bombs in some of the sealed rooms. The bombs are configured such that if even one photon of light is absorbed by the fuse (i.e. if someone looks into the room), the bomb will go off. Each room has an input and output port which can be hooked up to external devices. An empty room will let light go from the input to the output ports unaffected, whilst a room with a bomb will explode if light is shone into the input port and the bomb absorbs even just one photon. Your task is to find a way of determining whether a room has a bomb in it without blowing it up, so that specialised (limited and expensive) equipment can be devoted to defusing that particular room. Design a scheme that allows you – at



²This is a slightly modified version of a bomb testing problem described by Avshalom Elitzur and Lev Vaidman in *Quantum-mechanical interaction-free measurement*, *Found. Phys.* **47**, 987-997 (1993).

least part of the time – to decide whether a room has a bomb in it without blowing it up. If you iterate the procedure, what is its overall success rate for the detection of a bomb without blowing it up?

- (5) Assume that the two beam splitters in the interferometer are different. A beamsplitter which reflects incoming light with probability r and transmits with probability $1 - r$ is described by the matrix

$$B = \begin{bmatrix} \cos \theta & i \sin \theta \\ i \sin \theta & \cos \theta \end{bmatrix},$$

where r is parametrised as $\cos^2 \theta$. Let the probability of reflection in the first beamsplitter be r and in the second one $1 - r$. Would the new setup improve the overall success rate of the detection of a bomb without blowing it up?

- (6) There exists a scheme, involving many beamsplitters and something called “quantum Zeno effect”, such that the success rate for detecting a bomb without blowing it up approaches 100%. Try to work it out or find a solution on internet.
- (7) Show, by expressing matrices in terms of their components or otherwise, that that $(UV)^\dagger = V^\dagger U^\dagger$ and that the product of two unitary matrices is another unitary matrix. The set of all unitary $N \times N$ matrices with the matrix multiplication forms a non-Abelian group called $U(N)$.
- (8) Any unitary $N \times N$ matrix has N^2 complex entries, or, if we were to view each complex number as a pair of real numbers, $2N^2$ real entries. They are not independent from each other because of the unitarity condition. Show that any unitary $N \times N$ matrix can be parametrised by N^2 independent real parameters.
- (9) Check that any matrix of the form,

$$e^{i\phi} \begin{bmatrix} \cos \theta e^{i\alpha} & -\sin \theta e^{i\beta} \\ \sin \theta e^{-i\beta} & \cos \theta e^{-i\alpha} \end{bmatrix} \tag{15}$$

where α, β, θ and ϕ are independent real parameters, is unitary. In fact, any 2×2 unitary matrix can be written in this form.

- (10) The exponent of matrix A is defined as

$$e^A = \mathbb{1} + A + \frac{A^2}{2!} + \frac{A^3}{3!} + \dots = \sum_{n=0}^{\infty} \frac{(A)^n}{n!}$$

Show that if H is self-adjoint (Hermitian), that is $H = H^\dagger$, then $U = e^{iHt}$ is unitary for any real t . Many quantum evolutions are expressed in this way. This is because matrix H , known as the Hamiltonian, is related to energies, which are measurable physical quantities, and t stands for time.

- (11) Show that for any real α and for any A such that $A^2 = \mathbb{1}$

$$e^{i\alpha A} = \cos \alpha \mathbb{1} + i \sin \alpha A. \tag{16}$$

- (12) A qubit (spin one-half particle) initially in state $|0\rangle$ (spin up) is placed in a uniform magnetic field. The interaction between the field and the qubit is described by the Hamiltonian

$$H = \omega \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix},$$

where ω is proportional to the strength of the field. What is the state of the qubit after time $t = \pi/4\omega$?

Hint: Consider the setup where the input and output ports are hooked up in one of the arms of a Mach-Zehnder interferometer.

In Earth’s magnetic field, which is about 0.5 gauss, the value of ω is of the order of 10^6 cycles per second.